

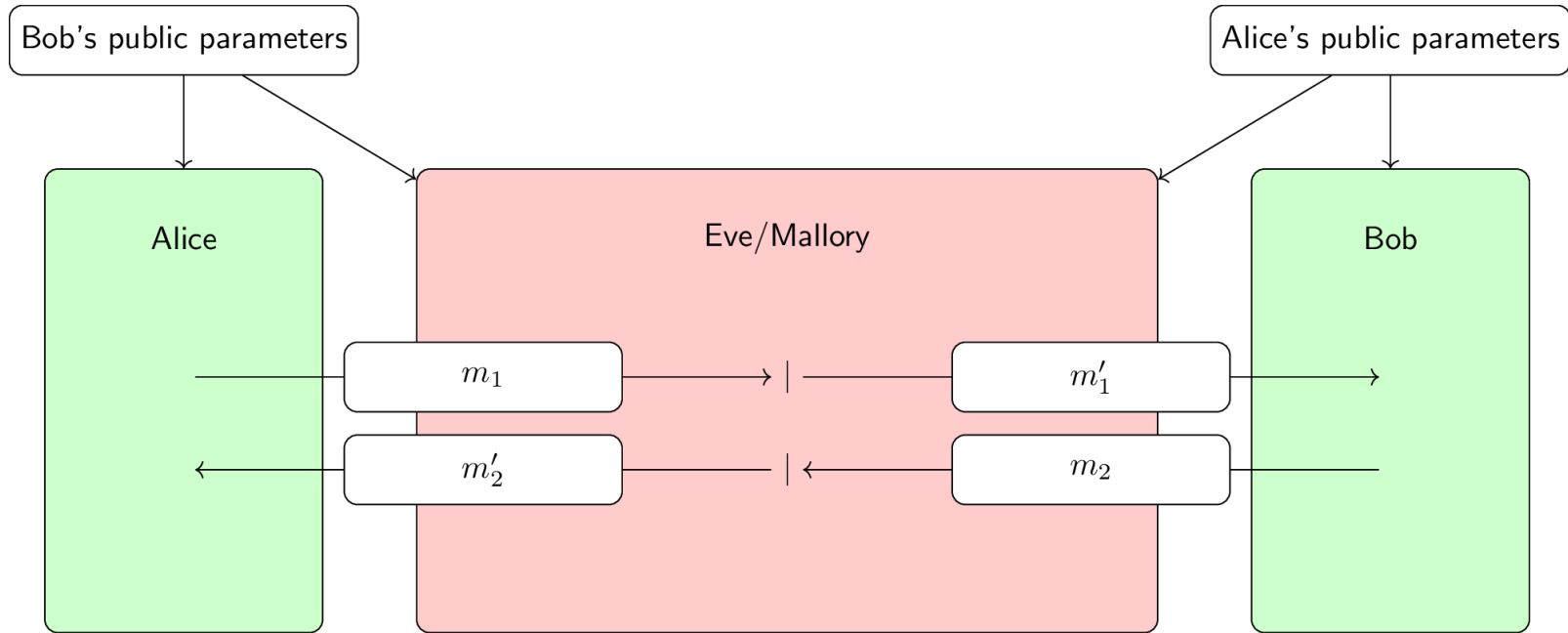
# Advanced Security for Systems Engineering – VO 10: Applied Cryptography

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# Threat Model: Passive vs Active

# Threat Model



Passive:  $m_i = m'_i$

- Depend on exact model
  - Passive: eavesdropping
  - Active: tampering with, blocking, delaying, reordering messages
  - Advanced active: corrupting some peers, etc (multiparty setting)
  
- Mostly: Probabilistic Polynomial Time (PPT) adversary
- If unsure, use most conservative model/most powerful adversary
- Always assume active adversary in a networking setting

- Ciphertext Indistinguishability
- Semantic Security
- Chosen Plaintext Attack
- Chosen Ciphertext Attack
- IND-CPA, IND-CCA, IND-CCA2

Blackboard



Brainstorming Attacks

# Common Attacks against Crypto

- Use of wrong protocol, insufficient security guarantees
- Protocol errors
- Implementation errors
- Side-channel attacks, Fault injection
- Statistical attacks, attacks on traffic patterns
- Compromise infrastructure, trust anchors

Which are Out-of-Model attacks?

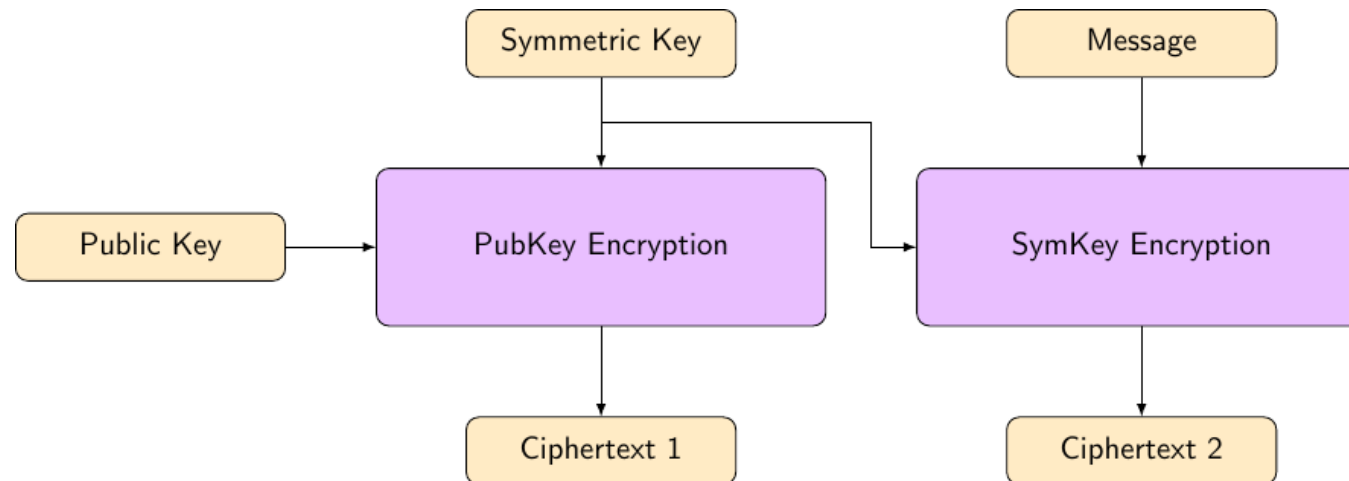


# Encryption Schemes

# Encryption Algorithms: Keywords

- Symmetric, Secret-key:  $m = D(k, E(k, m))$ 
  - 3DES, AES, (X)Salsa20, ChaCha
  - Fast, but Key Distribution problem
- Asymmetric, Public-key:  $m = D(sk, E(pk, m))$ 
  - RSA, ElGamal, Elliptic Curves

- Hybrid:



From Oneway Function/PRP to Secure Cryptographic Scheme

1. Oneway function (with trapdoor)/pseudorandom permutation (PRP)
2. Hardness assumptions
3. Threat model and goals (IND-CCA, IND-CPA)
4. Secure cryptographic scheme with reduction to hardness assumption

## ■ Assumption:

- Hardness related to Integer Factorization problem

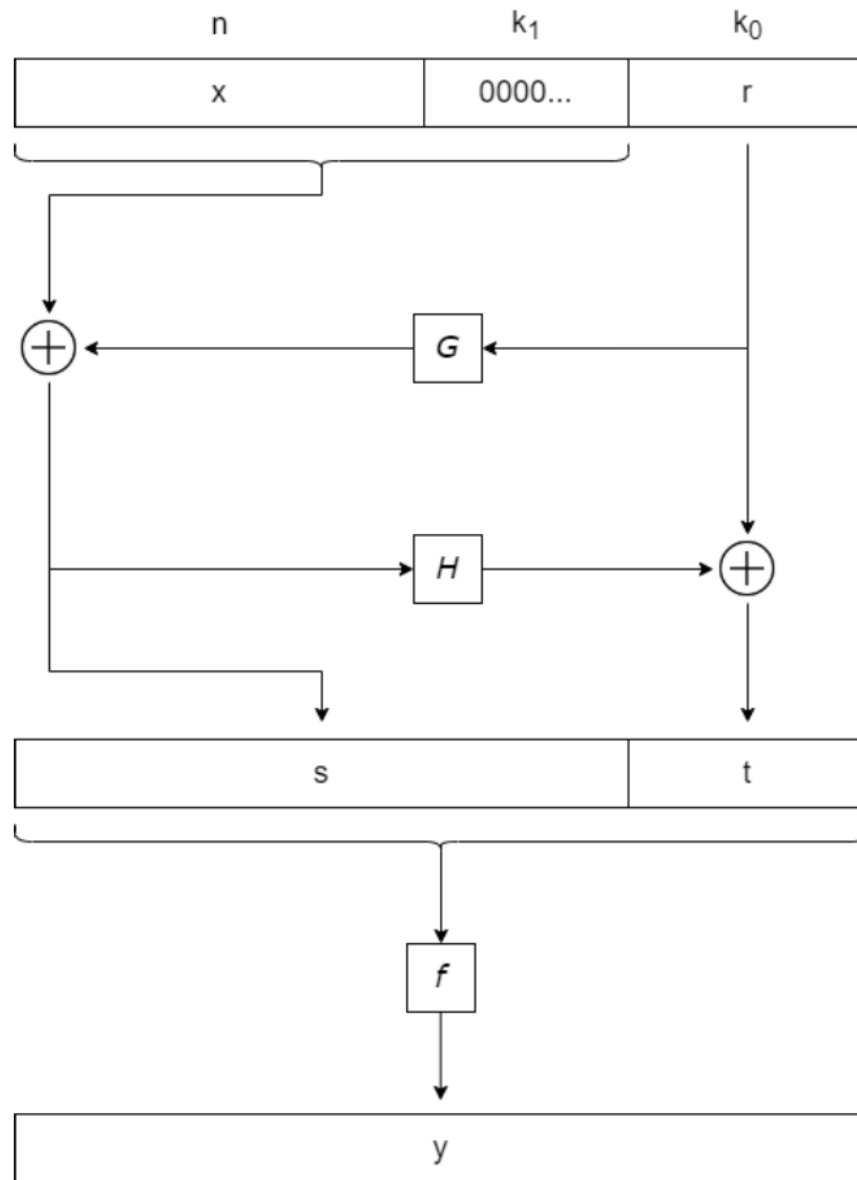
## ■ Basic Primitive:

- $N = p \cdot q$  with  $p, q \in \mathbb{P}$
- Operations are computed  $\text{mod } N$
- $sk : d$   $pk : e$  with  $e \cdot d = 1 \text{ mod } \phi(N)$
- $E : m^e$
- $D : m^d$

## ■ Secure Scheme:

- Never use plain (textbook) RSA, use OAEP or at least PKCSv1.5

# IND-CCA Security for RSA: OAEP



## ■ Assumption:

- Hardness of Discrete Logarithm, Decisional Diffie-Hellman (DDH)

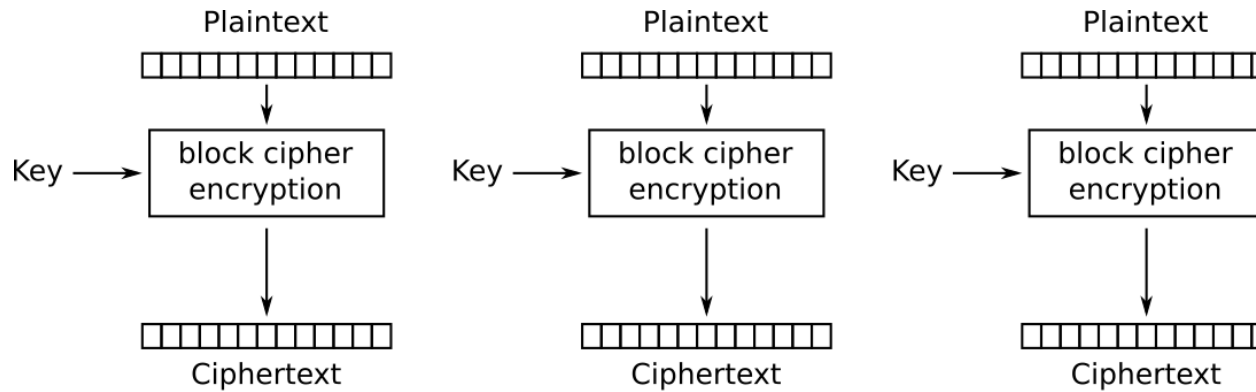
## ■ Basic Primitives (ElGamal)

- $p \in \mathbb{P}$ ,  $g$  is generator of  $\mathbb{Z}_p$
- Operations are computed  $\text{mod } P$
- $sk : x$   $pk : g^x$  with  $x$  uniform random sampled in  $\mathbb{Z}_p$
- $E : (c_0 = g^y, c_1 = pk^y \cdot m)$  with  $y$  uniform sampled in  $\mathbb{Z}_p$
- $D : \frac{c_1}{(c_0)^x}$

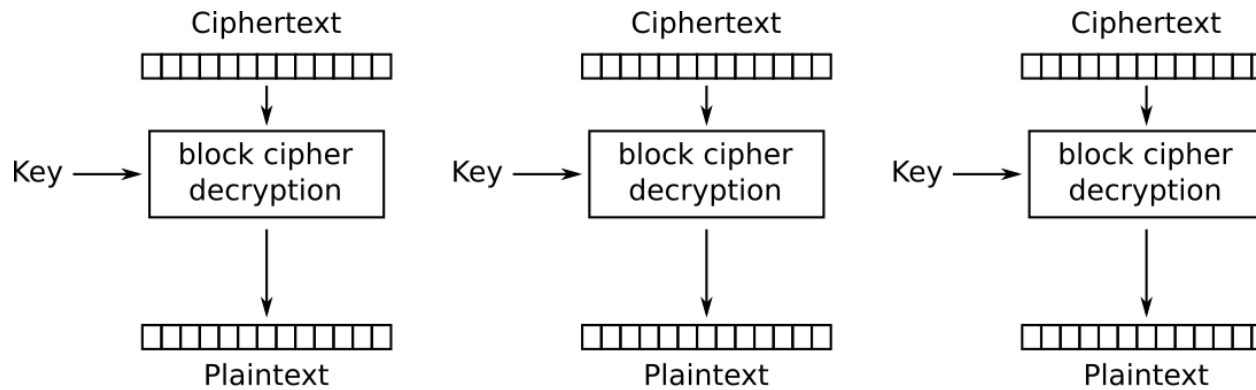
## ■ Secure Scheme:

- Cramer-Shoup extends Elgamal and is IND-CCA2 secure (DDH)

# Electronic Codebook (ECB) Mode



Electronic Codebook (ECB) mode encryption



Electronic Codebook (ECB) mode decryption

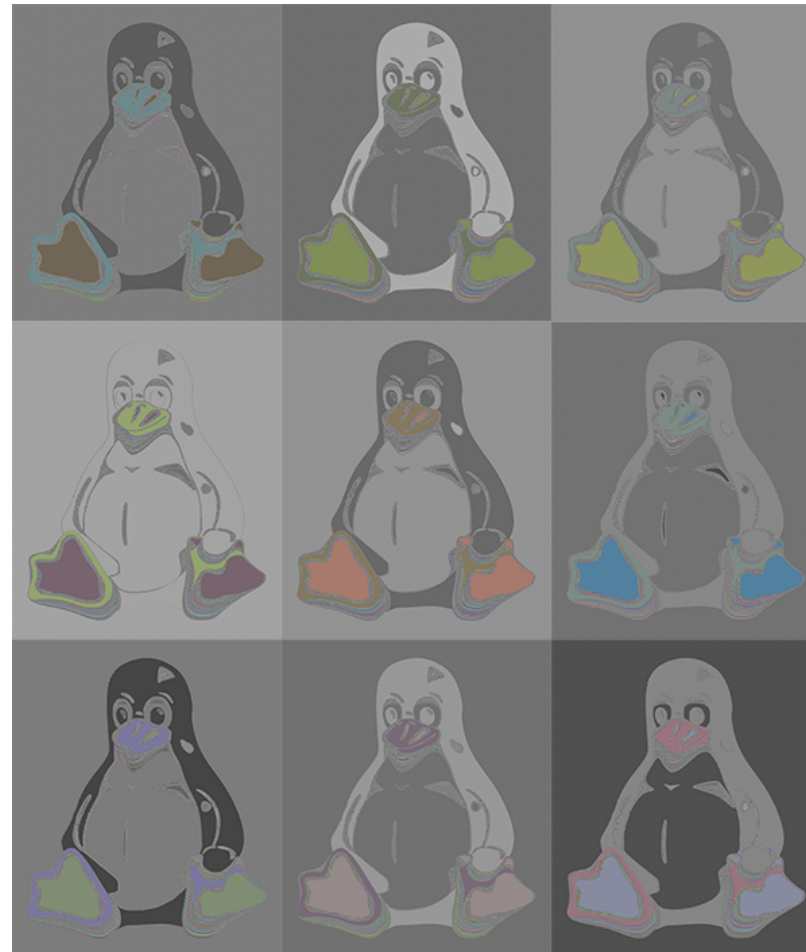
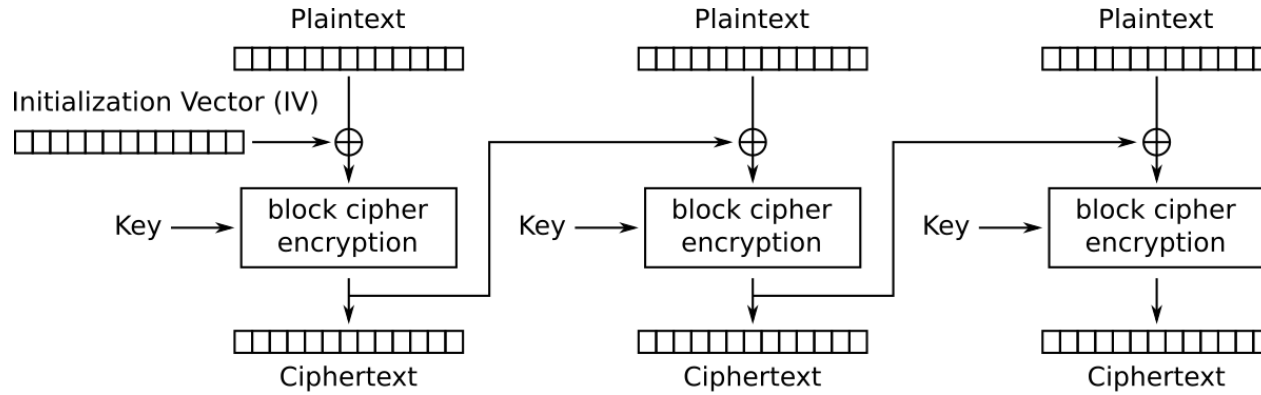
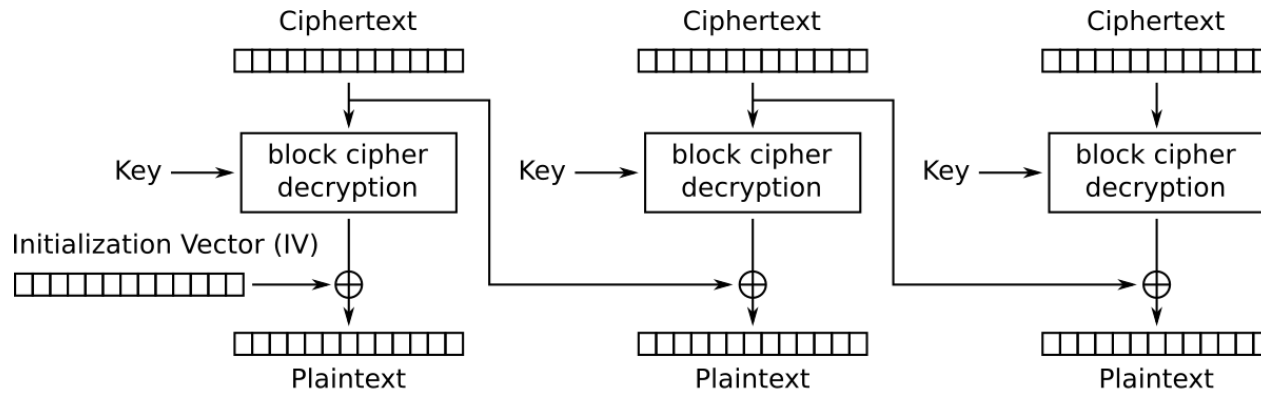


Figure 1: <https://blog.filippo.io/the-ecb-penguin/>





Cipher Block Chaining (CBC) mode encryption



Cipher Block Chaining (CBC) mode decryption

- Jonathan Katz, Yehuda Lindell: Introduction to Modern Cryptography, CRC Press, 2014
- Vaudenay: Security Flaws Induced by CBC Padding. Applications to SSL, IPSEC, WTLS. EUROCRYPT'02
- Boeck, et al: Return Of Bleichenbacher's Oracle Threat (ROBOT), Usenix Sec'18
- NaCl Library, [nacl.cr.yp.to](http://nacl.cr.yp.to)
- Libsodium Library, [libsodium.org](http://libsodium.org)
- [Boneh] Dan Boneh (Stanford): Online Cryptography Class. <http://crypto-class.org>

**Thank's for your attention!**

`https://security.inso.tuwien.ac.at/`

